

# Design of Low Power High Speed MultiplierUsing SPST Technique

C.Vinay<sup>1</sup>, K.Sai keerthi<sup>2</sup>, N.Rahul<sup>3</sup>, K.Praveen<sup>4</sup> <sup>123</sup>B.E student, <sup>4</sup>Assistant professor, ECE Dept. Department of Electronics and Communication Engineering Matrusri Engineering College, Saidabad, Hyderabad 500 059.

# Abstract

With the recent rapid advances in multimedia and communication systems, real-time signal processing like audio signal processing, video/image processing, or large-capacity data processing are increasingly being demanded. The multiplier is the essential elements of the digital signal processing such as filtering, convolution, and Inner products. The objective of a good multiplier is to provide a physically compact, good speed and low power consumption. In this project, we proposed a high speed and low-power multiplier by adopting the new SPST implementing approach. This multiplier is designed by equipping the Spurious Power Suppression Technique (SPST) on a modified Booth encoder which is controlled by a detection unit using an AND gate. The modified booth encoder will reduce the number of partial products generated by a factor of 2. The SPST adder will avoid the unwanted addition and thus minimize the switching power dissipation. In this project for simulation we use Model sim for logical verification, and further synthesizing it on Xilinx-ISE.

*Keywords:* multiplier, low power, spurious power suppression technique, modified booth encoder.

# I. INTRODUCTION:

Fast multipliers are the essential elements of the digital signal processing such as filtering, convolution. Most digital signal processing methods use nonlinear functions such as discrete cosine transform or discrete wavelet transform. Because they are basically accomplished by repetitive application of multiplication and addition, the speed of the multiplication and addition arithmetic's determines the execution speed and performance purpose processors today, especially since the media processing took off. In the past multiplication was generally implemented via a sequence of addition, subtraction, and shift operations. Multiplication can be considered as a series of repeated additions. The number to be added is the multiplicand, the number of times that it is added is the multiplier, and the result is the product. Each step of addition generates a partial product. In most computers, the operand usually contains the same number of bits. When the operands are interpreted as integers, the product is generally twice the length of operands in order to preserve the information content. It is possible to decompose multipliers into two parts. The first part is dedicated to the generation of partial products, and the second one collects and adds them. The basic multiplication principle is twofold i.e. evaluation of

partial products and accumulation of the shifted partial products. It is performed by the successive additions of the columns of the shifted partial product matrix. The 'multiplier' is successfully shifted and gates the appropriate bit of the 'multiplicand'. The delayed, gated instance of the multiplicand must all be in the same column of the shifted partial product matrix. They are then added to form the product bit for the particular form. Multiplication is therefore a multi operand operation. To extend the multiplication to both signed and unsigned numbers, a convenient number system would be the representation of numbers in two's complement format. Booth multiplication is a technique that allows for smaller, faster multiplication circuits, by recoding the numbers that are multiplied. It is the standard technique used in chip design, and provides significant improvements over the "long multiplication" technique. A standard approach that might be taken by a novice to perform multiplication is to "shift and add", or normal "long multiplication". That is, for each column in the multiplier, shift the multiplicand the appropriate number of columns and multiply it by the value of the digit in that column of the multiplier, to obtain a partial product. The partial products are then added to obtain the final result.

### 2. MODIFIED BOOTH ENCODER

In order to achieve high-speed, multiplication algorithms using parallel counters, such as the modified Booth algorithm has been proposed, and some multipliers based on the algorithms have been implemented for practical use. This type of multiplier operates much faster than an array multiplier for longer operands because its computation time is proportional to the logarithm of the word length of operands. Booth multiplication is a technique that usually allowed for smaller and faster multiplication circuits, by recoding the numbers that are multiplied. It is possible to reduce the number of partial products by half, by using the technique of radix-4 Booth recoding. The basic idea is that, instead of shifting and adding for every column of the multiplier term and multiplying by 1 or 0, we only take every second column, and multiply by  $\pm 1, \pm 2$ , or 0, to obtain the same results. The advantage of this method is the halving of the number of partial products. To Booth recode the multiplier term, we consider the bits in blocks of three, such that each block overlaps theprevious block by one bit. Grouping starts from the LSB, and the first block only uses two bits of the multiplier. Figure 3 shows the grouping of bits from the multiplier term for use in modified booth encoding.



Case (1):

Volume: 04 Issue: 06 | June -2020

ISSN: 2582-3930

l	0	1	0	1	1	0	1	0	1	0

### Figure 1. Grouping bits from the multiplier term

Each block is decoded to generate the correct partial product. The encoding of the multiplier *Y*, using the modified booth algorithm, generates the following five signed digits, -2, -1, 0, +1, +2. Each encoded digit in the multiplier performs a certain operation on the multiplicand, *X*, as illustrated in Table 1.

Table1
Booth encoding table

Block	Re-coded digit	Operation of X
000	0	0X
001	+1	+1X
010	+1	+1X
011	+2	+2X
100	-2	-2X
101	-1	-1X
110	-1	-1X
111	0	0X

#### **3. SPURIOUS POWER SUPRESSION TECHNIQUE**

Figure 2 shows the five cases of a 16-bit addition in which the spurious switching activities occur. The 1st case illustrates a transient state in which the spurious transitions of carry signals occur in the MSP though the final result of the MSP are unchanged. The 2nd and the 3rd cases describe the situations of one negative operand adding another positive operand without and with carry from LSP, respectively. Moreover, the 4th and the 5th cases respectively demonstrate the addition of two negative operands without and with carry-in from LSP. In such cases, the results of the MSP are predictable Therefore the computations in the MSP are useless and can be neglected. The data is separated into the Most Significant Part (MSP) and the Least Significant Part (LSP). To know whether the MSP affects the computation results or not. The detection logic unit is to detect the effective ranges of the inputs.

A MSP=A [15:8]; B MSP=B [15:8] (1	)
A AND = A [15].A [14]A [8](2)	
B AND = B [15].B [14]B [8](3)	
NOD = A[1E] + A[1A] + A[0]	()

 $\overline{A \ NOR \ = \ A[15] + A[14] + \dots \dots A[8]}.....(4)$  $\overline{B \ NOR \ = \ B[15] + B[14] + \dots \dots B[8]}.....(5)$ 

 $Close = \sim ((A_{AND} + A_{NOR}).(B_{AND} + B_{NOR})).....(6)$ 

We derive the Karnaugh maps which lead to the Boolean equations (7) and (8) for the Carr\_ctrl and the sign signals, respectively.

 $Carr\_ctrl = (C_{LSP} \oplus A_{AND} \oplus B_{AND}) (A_{AND} + A_{NOR})$  $(B_{AND} + B_{NOR}).....(7)$ Sign =  $\overline{CLSP}$ . (A<sub>AND</sub> + B<sub>AND</sub>) + C<sub>LSP</sub> A<sub>AND</sub> B<sub>AND</sub>.....(8)

(128)0000000010 000000 (-128)1111111110000000 +(64)000000001 000000 +(192)0000000011000000 0000000011 000000 00000000001000000 (192)(64) Case (2): Case (3)  $(A_{13} A_{14}...A_{0}) = (11...1), (B_{13} B_{14}...B_{0}) = (00...0), C_{7} = 0$  $(A_{15} A_{14}...A_{8}) = (11...1), (B_{15} B_{14}...B_{8}) = (00...0), C_{7} = 1$ (-196)1111111100111100 11111111111000011 (-61)0000000011001100 +(204)+(+51)000000000110011 (+8)(-10) 111111111111110110 Case (4):  $(A_{15} A_{14}...A_{3}) = (11...1), (B_{15} B_{14}...B_{3}) = (11...1), C_{2} = 0$ Case (5):  $(A_{15}\;A_{14}...A_8)=(11...1),\,(B_{15}\;B_{14}...B_8)=(11...1),\,C_7=1$ 11111111111000011 (-61)(-196)1111111100111100 +(-205)1111111100110011 11111111111001100 +(-52)1111110110 (-266)(-248)1111111100001000 sign carr-ctrl Figure 2. Transition cases in multimedia / DSP

# processing

# 4. PROPOSED SPURIOUS POWER SUPRESSION TECHNIQUE

The SPST uses a detection logic circuit to detect the effective data range of arithmetic units, e.g., adders or multipliers. When the portion of data does not affect the final computing results, the data controlling circuits of the SPST latch this portion to avoid useless data transitions occurring inside the arithmetic units. Besides, there is a data asserting control realized by using registers to further filter out the useless spurious signals of arithmetic unit every time when the latched portion is being turned on. This asserting control brings evident power reduction. Figure 3 shows the design of low power adder/subtractor with SPST. Adder /subtractoris divided into two parts, the most significant part (MSP) and the least significant part (LSP). The MSP of the original adder/subtractor is modified to include detection logic circuits, data controlling circuits, sign extension circuits, logics for calculating carry in and carry out signals. The most important part of this study is the design of the control signal asserting circuits, denoted as asserting circuits in Figure 3. Although this asserting circuit brings evident power reduction, it may induce additional delay. There are two implementing approaches for the control signal assertion circuits. The first implementing approach of control signal assertion circuit is using registers. This is illustrated in Figure 4. The three output signals of the



detection logic were close signal, Carr\_ctrl signal, sign signal. The three output signals the detection logic unit are given a certain amount of delay before they assert. The delay $\phi$ , used to assert the three output signals, must be set in a range of  $\psi < \phi < \Delta$ , where $\psi$  denotes the data transient period and  $\Delta$ Denotes the earliest required time of all the inputs. This will filter out the glitch signals as well as to keep the computation results correct. The restriction that  $\phi$  must be greater than  $\psi$  to guarantee the registers from latching the wrong values of control usually decreases the overall speed of the applied designs.







This issue should be noticed in high- end applications which demands both high speed and low power requirements. To solve this problem we adopt the other implementing approach of control signal assertion circuit using AND gate.



Figure 5.Detection Logic circuits using AND gate

## 5. PROPOSED LOW POWER HIGH SPEED MULTIPLIER

The proposed high speed low power multiplier is designed by equipping the SPST on a tree multiplier. There are two distinguishing design considerations in designing the proposed multiplier as listed in the following

A. Applying the SPST on the Modified Booth Encoder

Figure 6 shows a computing example of Booth multiplying two numbers "2AC9" and "006A". The shadow denotes that the numbers in this part of Booth multiplication are all zero so that this part of the computations be neglected. Saving can those computations can significantly reduce the power consumption caused by the transient signals. According to the analysis of the multiplication shown in figure 6, we propose the SPST-equipped modified-Booth encoder, which is controlled by a detection unit. The detection unit has one of the two operands as its input to decide whether the Booth encoder calculates redundant computations as shown in Figure 7. The latches can, respectively, freeze the inputs of MUX-4 to MUX-7 or only those of MUX-6 tMUX-7 when the PP4 to PP7 or the PP6 to PP7 are zero; to reduce the transition power dissipation. Figure 8. Shows the booth partial product generation circuit. It includes AND/OR/ EX-OR logic.



Volume: 04 Issue: 06 | June -2020

2AC9 0010101011001001	
006A x) 00000000011010100	
	DDO
11010101001101110	PP0
11101010100110111	PP1
11101010100110111	PP2
00101010110010010	PP3
00000000000000000	PP4
0000000000000000	PP5
0000000000000000	PP6
+)00000000000000000	PP7
0000000000100011011011100111010 (	11B73A





Figure 7.SPST Equipped Modified Booth Encoder



**Figure 8.Booth Partial Product Selector Logic** 

B. Applying the SPST on the Compression Tree The proposed SPST -equipped multiplier is illustrated in figure 9. The PP generator generates five candidates of the partial products, i.e., {-2A,-A, 0, A, 2A. These are then selected according to the Booth encoding results of the operand B. When the operand besides the Booth encoded one has a small absolute value, there are opportunities to reduce the spurious power dissipated in the compression tree.



**Figure 9.Proposed High Performance Low Power Equipped Multiplier** 

### 6. COMPARISION

The Proposed multiplier results are very good when compared with the existing methods and comparison table is given by following table.

Multiplier	CSA	SPST
type	MULTIPLIER	MULTIPLIER
vendor	Xilinx	Xilinx
family	Spartan 6	Spartan 6
Estimated	28.6ns	14.51ns
delay		
Power	0.218mw	0.082mw
dissipation		

### 7. RESULTS

In this project we are evaluating the performance of the proposed high speed low power multiplier by



Volume: 04 Issue: 06 | June -2020

using VHDL coding.Power report and delay report we are synthesizing these multipliers using Xilinx.

						1,831,2 <del>41</del> ps			
Name	Value	(1,831,240 ps	1,831,241.ps	1,831,242ps	1,831,243 ps	1,831,2 <del>41</del> ps	1,831,245 ps	1,831,246 ps	1,831,247 ps
dock	1					1			
Multiple(150)	006a					005a			
Multiplicand[15:0]	2809					28:5			
▶ 👹 PP_out_final1(16:0)	laafe	1				laafe			
▶ 👹 PP_out_final2(16:0)	14537					1:1537			
🕨 💐 PP_out_finaB(16:0)	14537					1:1537			
🕨 💘 PP_out_finaH(16:0)	05592					05592	2		
🕨 👹 PP_out_final5(16:0)	00000					00000			
🕨 💐 PP_out_final6(16:0)	00000					00000			
▶ 💘 PP_out_fina(7(16:0)	00000					00000			
▶ 💐 PP_out_finaB(16:0)	00000					00000			
🕨 🙀 Final_OUT(31:0)	0011b73a					0011b73a			
🕨 💐 stage1_add_out1[23x]	11b73a					11b73a			
▶ 💐 stage1_SPST_out1(19:0)	00100					00000			
🕨 👹 stagel_SPST_out2(19:0)	00000					00000	8		
🕨 💐 stage2_SPST_out1(23:0)	000000					000000			
Final_OUT_reg[31.0]	0011b73a					0011b73a			

Figure 10.Simulation results of proposed multiplier

A	8	C	D	E	F	G	Н	1	1	K	L	М	N
Device	-		On Chip	Power (W)	Used	Available	Utilization (%)		Supply	Summary	Total	Dynamic	Quiescent
Family	Artix7		Logic	0.000	855	63400	1		Source	Votage	Current (A)	Current (A)	Current (A)
Part	xc7a100t		Signals	0.000	1089	-	-		Vocint	1.000	0.017	0.000	0.01
Package	csg324		10s	0.000	65	210	31		Vccaux	1.800	0.013	0.000	0.01
Temp Grade	Commercial	¥	Leakage.	0.082		(	· · · · · ·		Vcco18	1.800	0.004	0.000	0.00
Process	Typical	¥	Total	0.082					Vccbram	1.000	0.000	0.000	0.00
Speed Grade	-3					_			Vccadc	1.710	0.020	0.000	0.02
					Effective TJA	Max Ambient	Junction Temp						
Environment			Therma	Properties	(C/W)	(0)					Total	Dynamic	Quiescent
Ambient Temp (C)	25.0				4.6	84.6	25.4		Supply	Power (W)	0.082	0.000	0.08
Use custom TJA?	No												
Custom TJA (C/W)	NA												
Aitflow (LFM)	250	•											
Heat Sink	Medium Profile	¥											
Custom TSA (C/W)	NA												
Board Selection	Medium (10"x10")												
# of Board Layers	12 to 15												
Custom TJB (C/W)	NA												
Board Temperature (	o NA												
Characterization													
Production	v1.0,2012-07-11												

Figure 11. Power report of proposed multiplier

MITVOV (CT=>0	1	0 023	0.000	SD2/ADD MED2/n5/Madd n0004 Madd ouc25 (SD2/ADD MSD2/n5/Madd n0004 Madd ouc25)
MITVOV (PT-SO	1	0 022	0.000	STS/ADD HOTS/HS/HS/AS BRARA MS/AS ADD S (STS/ADD HOTS/HS/HS/A HD/A HS/AS ADD S)
MODOW OT NO		0.023	0.000	arz/ADD Marz/Ba/Medd BUUUN MEDA CYSAZ (arz/ADD Marz/Ba/Medd BUUUN MEDA CYSAZ)
KURCTICI-20	-	0.379	0.073	SP2/ADD MSP2/PS/Madd R0004 Madd XGEV47 (SP2/ADD MSP2/PSeudo SumV47)
L0T6:11=>0	1	0.097	0.024	SPA/ADD MSP1/N1/H NOKZ SWU (NB)
LUT6:10=>0	2	0.097	0.516	SP3/ADD_MSP1/n1/B_nor2 (SP3/ADD_MSP1/n1/B_nor2)
LUT6:13->0	19	0.097	0.767	SP3/ADD_MSP1/n2/Mmux_closel_(SP3/ADD_MSP1/MSP_A_Out<0> mand)
LUT6:11->0	1	0.097	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd lut<0> (SP3/ADD MSP1/n5/Madd n0004 Madd lut<0>
MEXCY:S=>0	1	0.353	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd cy<0> (SP3/ADD MSP1/n5/Madd n0004 Madd cy<0>)
MIXCY:CI->0	1	0.023	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd cy<1> (SP3/ADD MSP1/n5/Madd n0004 Madd cy<1>)
MIXCY (CI->0	1	0.023	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd cv<2> (SP3/ADD MSP1/n5/Madd n0004 Madd cv<2>)
MUXCY:CI=>0	1	0.023	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd ev<3> (SP3/ADD MSP1/n5/Madd n0004 Madd ev<3>)
MIXCY:CI->0	1	0.023	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd evc4> (SP3/ADD MSP1/n5/Madd n0004 Madd evc4>)
MIXCY+CT=>0	1	0.023	0.000	SP3/ADD MSP1/n5/Madd n0004 Madd eve5> (SP3/ADD MSP1/n5/Madd n0004 Madd eve5>)
YORCY (CT=>0	î	0.370	0.556	SP1/ADD MSP1/n5/Madd n0004 Madd worc65 (SP1/ADD MSP1/Pnoudo Sunc65)
1075:11-20	1	0.097	0.000	ADDA/Madd OTT lute305 (ADDA/Madd OTT lute305)
MITVOV 8-20	0	0.357	0.000	ADDA (Model OPT THE SHOP (ADDA (Model OPT THE SHOP)
VODOV OT-SO	1	6 176	0.576	ADDA Manda (MTH Marchis (Dins) (MTH Mar 11 (DETD)
ADRUT ILI-AU	*	9.379	H.4.19	ADDA/PARKE ONE KOEKSEP (FIRME ONE ENG SECONF)
OBILE: 1=30		0.000		Pinal OUT reg st obor (Pinal OUT reg(s12)
	********			
Total		14.514n	1 (6.254	hs logic, K.260hs route)
			(43.11	logic, 56.9% route)

Figure 12. Delay analysis of proposed multiplie

### 8. CONCLUSION

In this project, we propose a high speed lowpower multiplier adopting the new SPST implementing approach. The proposed multiplier isdesigned by equipping the Spurious Power Suppression Technique (SPST) on a modified Booth encoder which is controlled by a detection unit using an AND gate. The modified booth encoder will reduce the number of partial products to half .TheSPST adder will avoid the unwanted addition and thus minimize the switching power dissipation. The SPST multiplier implementation using a detection logic uses an AND gates have an extremely high flexibility on adjusting the data asserting time. This facilitates the robustness of SPST can attain speed improvement and power reduction in the Modified booth encoder. This design can be verified using Modelsim and Xilinx using Verilog.

# 9. REFERENCES

[1] Sathya, A.; Fathimabee, S.; Divya, S., "Parallel multiplier-accumulator based on radix-2 modified Booth algorithm by using a VLSI architecture," Electronics and Communication Systems (ICECS), 2014 International Conference on , vol., no., pp.1,5, 13-14 Feb. 2014

[2] Suriya, T.S.U.; Rani, A.A., "Low power analysis of MAC using modified booth algorithm," Computing, Communications Networking and Technologies (ICCCNT),2013 Fourth International Conference on, vol., no., pp.1,5, 4-6 July 2013

[3] Naveen Kumar, Manu Bansal, AmandeepKaur, "Speed Power and Area Efficient VLSI Architectures of Multiplier and Accumulator," International Journal of Scientific & Engineering Research Volume 4, Issue 1, January-2013

[4] J. Y. YahwanthBabu, P. Dinesh Kumar, "A new multiplier Accumulator architecture based on high accuracy modified Booth algorithm," International journal of advanced research in computer Engineering & Technology, vol.2,no.3,pp.1036-1040,Mar 2013.



Vol. 3 No. 4 April,2012
[6] B. Jyothirmai, M. Premalatha, "Design and implementation of parallel MAC by Booth algorithm," International Journals of scientific Research,vol.1,no.6, pp.78-79, Nov.2012.

[7] AdalankiPurna Ramesh, Dr. A. V. N. Tilak, Dr. A. M. Prasad, "Efficient implementation of 16 bit multiplier, Accumulator using Radix – 2 modified Boothalgorithm and SPST adder using Verilog," International journal of VLSI design and communication systems, vol. 3, no.3, pp.107-118, Jun. 2012.

[8] PratibhadeviTapashetti, Dr. Praveen Kumar, "Compatible Architecture of MAC, Based on Modified Booth Algorithm". International Journal of Emerging Technology and Advanced Engineering. Oct 2017.

[9] Harilal, Durga Prasad, "High speed arithmetic Architecture of parallel multiplier Accumulator based on Radix-2 modified Booth algorithm," International journals of computational Engineering research, vol.2, no.8, pp.28-38, Dec.2012.

[10] V.ThirumalaRao, S. Girish Gandhi, C.Leela Mohan,
"Performance Evaluation of Parallel Multipliers for High Speed MAC Design", International Journal of Innovations in Engineering and Technology (IJIET). ISSN:2319-1058.
[11] Sareddy. Sindhuja Reddy, CH. BhanuPrakash.
"Design of a Novel Multiplier and Accumulator using Modified Booth Algorithm with Parallel Self-Time

Adder". International Research Journal of Engineering and Technology (IRJET) e-ISSN: 2395-0056 [12] A. R. Cooper, "Parallel architecture modified Booth

multiplier," Proc.Inst. Electr. Eng. G, vol. 135, pp. 125– 128, 1988

[13] Young Ho seo, Dong wookkim, "A new VLSI Architecture of parallel multiplier Accumulator based on Radix 2 Modified Booth algorithm," IEEE transactions on very large scale integration (VLSI) systems, vol. 18, no.2, pp. 201-208, Feb.2010.

[14] P. SasiBala& S. Raghavendra, "A New VLSI Architecture of Parallel Multiplier –Accumulator Based on Radix-2 Modified Booth Algorithm".International Journal of Instrumentation, Control and Automation (IJICA) ISSN: 2231-1890 Volume-1, Issue-2, 2011.

[15]. Pettam Ramesh, K. Uma Maheshwari, P.N, V. PurnaChanderRao. "A New VLSI Architecture Design based on Radix-2 Modified Booth Algorithm using ParallelMultiplier–Accumulator", International Journal of Scientific Engineering and Technology Research, 2319-8885, Vol.02, Issue.13, October-2013, Pages:1406-1415.

[16] S. Waser and M. J. Flynn, Introduction to Arithmetic for Digital Systems Designers. "New York: Holt, Rinehart and Winston, 1982."