

# DATA LEAKAGE AND PREVENTION SYSTEM IN SECURE CLOUD STORAGE USING LOCATION DEFINED NETWORK

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Abstract - Data sharing and access are capabilities businesses and organizations require the most these days. Remote working and mobile access to resources and collaboration platforms made it easier to access data and resources from anywhere, anytime. Employees want to access documents and email from different devices, and from various locations at a time. Access from untrusted networks is always a threat to businesses. This might result in data loss and overexposure of critical data. To mitigate the deficiencies of logical security mechanisms ,and coinciding with the trend of cyber-physical systems, security mechanisms have been proposed that integrate with the physical environment. To ensure that business's data and resources are safe. In this project we propose an innovative CodeFence that uses a location data and geospatial intelligence. Geospatial data analysis enhances understanding, insight, decision-making, and prediction. Location intelligence (LI) is achieved via visualization and analysis of geospatial data. Then we improve the security of data access in Data Server for a company or any other specific locations using the locationbased cryptosystem. CodeFence provides a means to secure sensitive information within an organization. It can be set to Off, On, Restricted View or Read Only. Once a geo-fenced boundary is defined, the opportunities what businesses can do is limited by only their creativity. The main benefit of setting up such a geo fence is in avoiding data leakage. Once defined the trusted network locations, no one can access data from a different network location/device. The experiment shows that our scheme is feasible in practical applications.

*Key Words*: geospatial, data access, CodeFence, Location, Security, Cloud Storage.

## **1.INTRODUCTION**

In today's fast-paced and interconnected world, data sharing and access have become critical capabilities for businesses and organizations to thrive. However, with remote working and mobile access to resources come the risks of accessing data from untrusted networks and the potential for data loss or overexposure. To address these challenges, security mechanisms have been proposed that integrate with the physical environment, coinciding with the trend of cyber-physical systems. In this context, we propose an innovative CodeFence that leverages location data and geospatial intelligence to enhance the security of data access in a company or any other specific location. By using geospatial data analysis and location intelligence, CodeFence provides a means to secure sensitive information within an organization, with the ability to be set to Off, On, Restricted View, or Read Only. This solution allows businesses to limit the opportunities for data leakage and restrict access to data only from trusted network locations. Our experiment shows that this scheme is feasible in practical applications, and we believe that it has the potential to be an effective solution for enhancing data security in today's business environment.

## 2. System Analysis 2.1. Existing System

Spatial-Temporal provenance Assurance with Mutual Proofs (STAMP) scheme. STAMP is designed for ad-hoc mobile users generating location proofs for each other in a distributed setting. However, it can easily accommodate trusted mobile users and wireless access points. STAMP ensures the integrity and nontransferability of the location proofs and protects users' privacy. A semi-trusted Certification Authority is used to distribute cryptographic keys as well as guard users against collusion by a light-weight entropy-based trust evaluation approach.

### xAd:

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This tool eliminates every form of assumption in marketing, because it serves messages based on your potential consumer's location. xAd has a proprietary platform that automatically creates boundaries around places often visited by a consumer. For example, a Restaurant, shopping mall. It's with these insights that marketers can target ads to their customers when they're within those locations.

### Koupon Media:

This tool prompts a targeted offer to shoppers when they're within the store. Koupon Media has features that study the behavioral attribute of buyers within the geofenced locations, and uses it to present the buyer with offers they can't resist while they are shopping.

### NinthDecimal:

This helps marketers to target consumers near their own stores or competitor's locations, with tangible media ads



through phone calls, appointment requests, and couponing.

### • Wal-Mart

It is another brand making it real big with geofencing. Their app comes with a store mode that picks up signals when a buyer is within the store, and delivers coupons and e-receipts.

## 2.2. Proposed System

This project will provide an introduction to Geo Server own authentication and authorization subsystems. such as from basic/digest authentication and CAS support, check through the various identity providers, such as Geo fence boundaries, MAC (Media Access Control), IP (Internet Protocol), as well as providing examples of custom authentication plug-in for Geo Server, integrating it in a home-grown security architecture. This system creates the victim file for wipe out the data, when the data is attempted to open outside of the geo fence.

## 2.2.1. Virtual Fence

The project propose а Geo-fencing (geofencing) is a feature in a software program that uses the global positioning system (GPS) to define geographical boundaries. To check whether a person is within a geofence range we can make use of different algorithms such as Ray-casting, Winding Number, TWC (Triangle Weight Characterization) and Circular Geofencing using Haversine Formula. Geofencing is security, when anyone enters or leaves a particular area, an alert passes to server. This system creates the victim file for wipe out the data, when the data is attempted to open outside of the geo fence.

## 2.2.2. Geospatial Intelligence Technology

Geo-fencing (geofencing) is a feature in a software program that uses the global positioning system (GPS) or radio frequency identification (RFID) to define geographical boundaries'-fencing allow an administrator to set up triggers so when a device enters (or exits) the boundaries defined by the administrator, an alert is issued. Geofence virtual barriers can be active or passive. Active geofences require an end user to opt-in to location services and a mobile app to be open. Passive geofences are always on; they rely on Wi-Fi and cellular data instead of GPS or RFID and work in the background. Geofences can be set up on mobile, tablet, and even desktop devices anywhere in the world.

## 3. System Architecture









## Level-1







# 4. Algorithms

# 4.1. Point in Geofence Framework and Algorithm:

### <u>Algorithm:</u>

Input: **r** is the radius of the geofence.  $\mathbf{g} = [g_0, g_i]$ g<sub>i</sub> is the position of the interest, go is the position of user. Output: true if **r** does not violate **g**, otherwise false 1: if pointInGeofence (g<sub>i</sub>, r) then 2: returntrue 3: end if 4: for all g<sub>0</sub>(i) in g<sub>0</sub> do 5: if pointInGeofence(r, g<sub>0</sub>(i)) then 6: returnfalse 7: end if 8: end for

9: return true

The above algorithm consists of the input parameters r and g. r = (x, y) is the current position to check for geofence violation. The geofence is specified by g = [gi, go] where gi is the keep-in geofence boundary polygon and  $G_0 = \{go1, ..., gOn\}$  is the set of keep-out boundaries.  $go_j$  is the j<sup>th</sup> of n keep-out geofence boundary polygons. The PointInGeofence () function can be implemented by any of the algorithm like Ray Casting, Winding Number, TWC and Circular Geofencing using Haversine formula

## 4.2. Ray Casting:

The Ray Casting algorithm determines whether or not the position of interest, Gi, is inside a given polygon p, by projecting an infinite ray from Gi. If the infinite ray intersects an odd number of polygon edges, then r is contained in p, otherwise, r is outside of p. As the Ray Casting algorithm iterates over all edges of p and does not have an initialization step, if the geofence boundaries change from one time step to the next, code execution and results of the Ray Casting algorithm are not impacted.

Algorithm:

Input: p is a simple polygon Gi is the position of interest buf is a buffer distance.

Output: true if p contains G<sub>i</sub>, otherwise false

1:  $\operatorname{count} = 0$ 

2: s is an infinite ray in the +y direction, originating atGi

3: for all edges e in p do

4: if  $\mathbf{\tilde{G}_{i}} \times \mathbf{\tilde{s}}$  within buf of ex then

- 5:  $e_{x,buf} = e_{x-2} * buf$
- 6: else

7: ebuf =e

8: end if

9: if  $G_i$  is within buf of e or  $e_{buf}$  then



## 4.3. Winding Number:

The winding number accurately determines if a point is inside a non-simple closed polygon. It does this by computing how many times the polygon winds around the point. The point is outside only when this "winding number" wn = 0; otherwise, the point is inside.

### Algorithm:

 $G_i = a point of Interest,$ V [] = vertex points of a polygon V[n+1] with V[n] = V[0]n = number of vertices Output: winding\_number (when the winding\_number =0, P is outside and winding\_number is non-zero if P is inside) PIP\_windingNumber (Point Gi, Point V [], int n) { Wn=0: 1: int *//* the winding number counter 2: for each edge E[j]: V[j] V[j+1] of the Polygon Do 3: if (E[j] crosses upward) if (P is strictly left ofE[j]) 4: 5: ++winding\_number; 6: end if 7: else if (E[j] crosses downward) 8: if (P is strictly right ofE[j]) 9: --winding\_number; 10: end if end if 11: 12: end for 13: return winding\_number;}

## 4.3. Triangle Weight Characterization

## (**TWC**):

Triangle Weight Characterization, consists of an initialization step and a run-time step as shown in Algorithm of TWC [4]. The initialization step must be executed for all keep-in and keep out geofences when the system first activates. If there are any changes to any of the geofence boundaries after the original initialization, each keep-in or keep-out geofence that is changed must be initialized again [4].

The initialization step subdivides each of the original geofences from simple polygons to y- monotone polygons and then to triangles [4]. The run-time step checks whether the position of interest is within each triangle. If the position of interest is inside any of the triangles, then it is within that polygon. Otherwise, it is outside the polygon [4]. Algorithm:

Input:

p is a simple polygon Gi is the position of interest Output: true if p contains  $G_i$ , otherwise false Initialization: 1: Divide p to m y-monotone polygons 2: for all y-monotone polygons M in p do 3: Divide polygon M to n triangles 4: end for Run-Time: 5: for all N triangles in p do 6: if N contains  $G_i$  then 7: return true 8: end if 9: end for

10: return false

# 4.4. Circular Geofencing Using Haversine

## Formula :

In the below algorithm, geofence of radius 'r' is created around the point Go The distance between Go and Gi is calculated using Haversine formula. The Haversine formula determines the great-circle distance between two points on a sphere given their longitudes and latitudes [6].

## Algorithm:

Input:Gi the position of the interest. Gi = [lati, longi] Go is the current position of the user. Go [ lato, longo] r is the radius of the geofence. Output: True if the  $G_i$  is within Geofence range of  $G_0$ checkWithInGeofenceRange () { 1: Distance = 02: Distance = getDistanceFrom Location (Go. Gi); 3: If (d < r)4: Return true 5: Else 6: Return false 7: End if } Get DistanceFromLoaction (Go, Gi) 8: radius =6371: //radiusofearth 9: dlat = deg2rad (lati, -lato); 10: dlong = deg2rad (longi, -long<sub>0</sub>); 11: a = Math.sin (dlat / 2) \* Math. sin (dlat/2)+ Math. cos (deg2rad(lato)) \* Math. cos (deg2rad(lati)) \* Math. sin (dlong/2) \*Math. sin (dlong/2); 12: c= 2 \* Math.atan2 (math. sqrt (a), Math.sqrt (1-a)); 13:  $D = R^*c; 14$ : Returnd; }



## **3. CONCLUSIONS**

In this project, we introduced a novel location-aware framework for provide data security, which enables the participation of workers without compromising their location privacy. We identified geo fencing as a needed step to ensure that data privacy is protected prior to workers consenting to a task. We provided heuristics and optimizations for determining effective geo fencing regions that achieve high task assignment rate with low overhead. It also generates the victim files; it automatically checks the geo - fencing boundary values and wipeout the system and files when geo fencing and MAC Address is mismatch.

## **3.1 Future Scope**

In the future, we plan to take into account more complicated policies to capture other privacy requirements other than the location. And also we insist this method to popular E-Mail Service Provider too.

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