

# Design of Energy Efficient Arithmetic VLSI Circuits

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**Abstract:** Advanced electronics have become an integral part of human lives, from Smartphones in our pockets to the modern vehicles on the road and to the satellites that orbit the earth. Power consumption has become a major constraint in these applications, often employing data-intensive digital signal processing systems and architectures. While enormous transistor density at the nanoscale has been increased in the Very Large-Scale Integration (VLSI) multicore era comparable to Moore's law, improving the performance of computing systems has become extremely difficult. Sustaining technological advancements for such energy-constrained devices has paved the way for new areas in research.

“Approximate Computing (AC)” has become one of the most promising paradigms, a field that has gained significant attention in the quest for energy efficiency in recent years. AC algorithms are numerically approximate rather than accurate. Inherent error resilience is the primary source of motivation behind AC. Recent research demonstrates that error resilience is pervasive in cyber search, deep learning, multimedia, recognition and data mining. Arithmetic units are the fundamental building blocks in most of the data dominated applications at the micro-architectural level of abstraction. AC is optimal for power and area efficient arithmetic circuits.

## 1. Introduction

The nanoscale miniaturization of modern electronics has resulted in numerous consumer devices spanning a broad spectrum of applications. High performance digital signal processors with embedded integrated circuits are in great demand for portable mobile devices. The steadfast expansion of big data processing, Internet of Things (IoT) and machine learning in addition to Multimedia necessitates a huge amount of data and massive computations at lower power budgets. Energy efficient arithmetic circuits are imperative requirement, that form the fundamental of digital systems.

Figure 1.1 shows two tasks that deal with division between 500 and 21. When asked whether the result of division between 500 and 21 is greater than 10, the human brain would right away respond "Yes," Task division is done but not to the fullest accuracy. However, if the question is whether the result is greater than 23.8, we would likely take a little longer and work even harder. However, system hardware and software typically do computations with consistent accuracy levels every time.

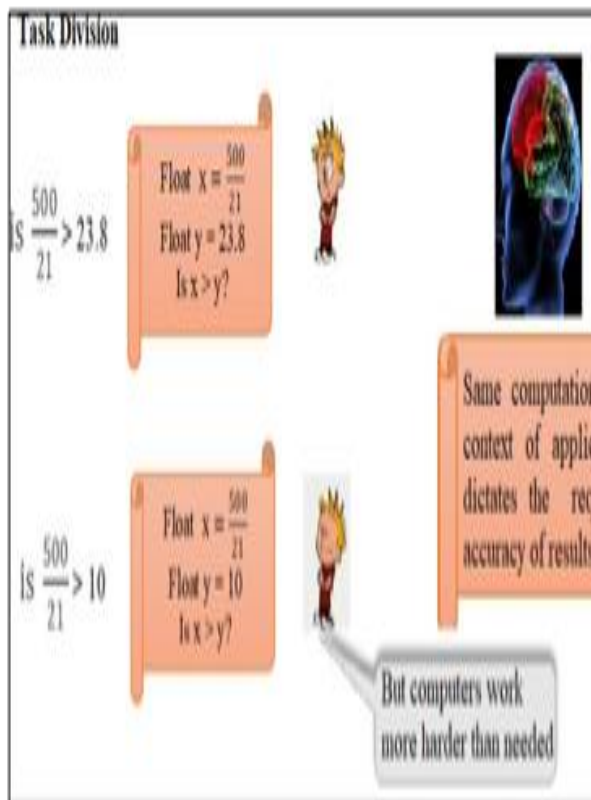


Fig 1.1 Task division

In a larger sense, edge devices coupled with battery power driven fog nodes desire minimal power. This demand emerges as we are swiftly approaching an era dominated by a society built on the efficiency across several error-resilient IoT and cyber-physical systems with cloud factors, various technological requirements infrastructure as its foundation. Facilitating and emphasis on green computing, efficient communication between these devices becomes paramount to reduce network latency and energy usage while ensuring that remote processing is more feasible.

## 2. APPROXIMATE COMPUTING(AC)

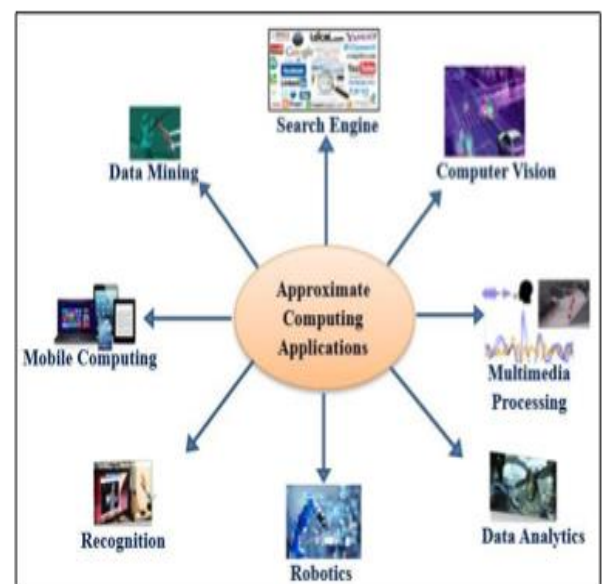
Digital Signal Processing (DSP) blocks form the core of various multimedia applications in which the ultimate result is either an image or

video targeting people's needs. However, humans have limited perceptual abilities when interpreting a picture or video. Inheriting the advantage, it is possible for the outputs of the existing digital logic circuits to be approximate with a trivial trade-off in accuracy. This field of research is generally known as approximate or error-tolerant computing. AC depends on the ability of computational systems to accept some inaccuracies being executed using approximate designs.

AC primarily depends on inherent error resiliency. Inherent error resilience or tolerance is the property of an application to generate acceptable outputs instead of being accurate, even though some aspects of its underlying computations are approximate or imprecise. Recent research demonstrates that a wide range of applications ( $\approx 83\%$ ) are inherently error-tolerant. AC is well-suited for error-tolerant multimedia applications.

## 3. NECESSITY FOR APPROXIMATE COMPUTING

The following subsections concisely



## Fig 2. Approximate computing in multimedia application

Noisy and redundant input data: Many DSP systems receive noisy inputs from the outside world. Applications with resilience to noisy inputs also have tolerance to approximations in computations.

- **Absence of a Unique Golden Output:** In some scenarios, such as those involving web search engines direct answers are not possible, but multiple valid responses are acceptable. A noteworthy example of error resilience is frequently cited as the vast family of Recognition, Mining and Synthesis (RMS) applications. Slightly modifying the outputs of internal operations may not degrade the end results.
- **Limited Human Perception:** Due to limitations in human perception, these errors are difficult to distinguish in applications like image, audio and video processing. The majority of the applications that belong to this category are found in the domains of multimedia and telecommunication

- **Self-healing:** In certain applications involving iterative refinement, errors due to approximations in one iteration can potentially get healed/recovered subsequently. Certain computational patterns favor the mitigation or the rejection of errors.

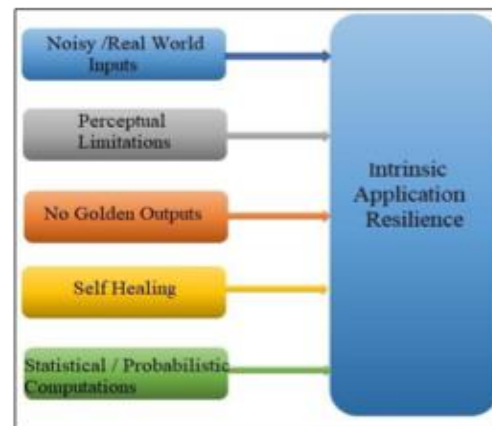


Fig 3. Various sources of inherent resilience

- **Statistical / Probabilistic computation:** In this approach, algorithms and techniques are designed to handle uncertainty and variability in data. These applications employ statistical computations that attenuate or cancel errors. This method helps in enhancing the reliability and accuracy of computations.

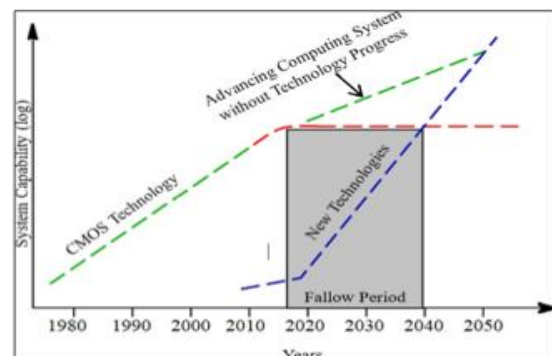
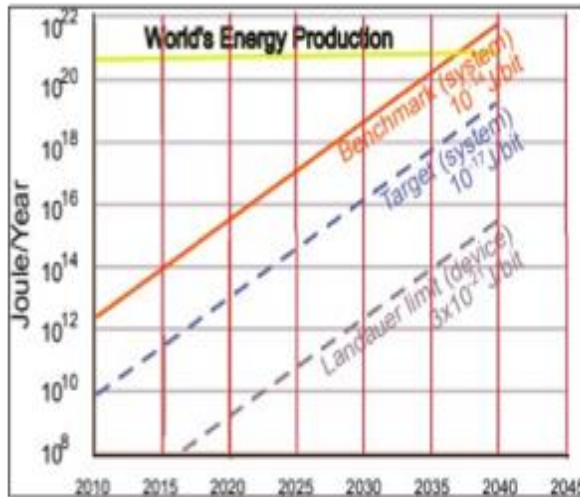


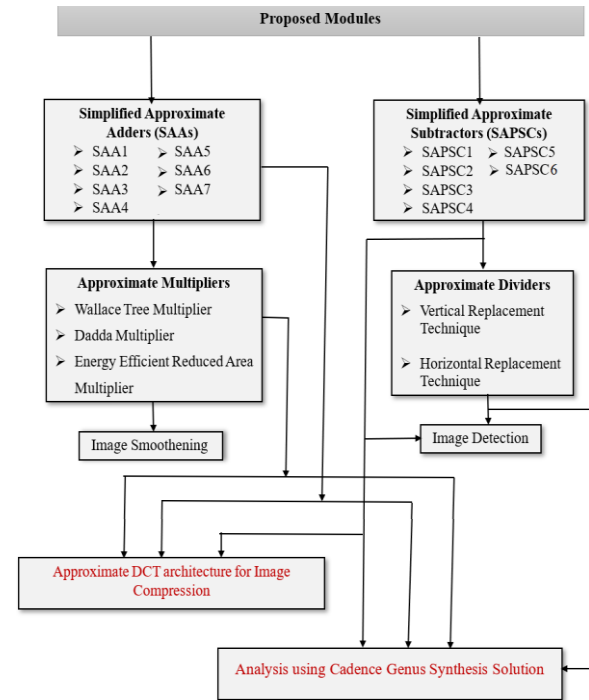
Fig 4. Harvest new gains in the "Fallow Period"



**Fig 5. Progress in Energy consumption vs. the Global energy production**

Approximate computing intentionally introduces trivial errors by emphasizing efficient enhancements in power, delay and circuit complexity across wide range of applications. Data processing has significant challenges in producing the perfect output. For a ‘good enough’ output that can be acceptable by the end user, accuracy can be condensed at various abstraction levels in the processing circuits.

The effectiveness of AC techniques in the embedded Support Vector Machine, Tensorial kernel circuit in tactile sensing systems is demonstrated by implementing approximate multipliers in the Singular Value Decomposition. AC using an end-to-end Trainable neural Network (AXNet) is suggested in combining dual NNs such as an approximator and a predictor into a comprehensive end-to-end trainable Neural Network (NN) to ensure the approximation quality by using the multi-task learning concept.



**Fig 6. Overview of the proposed modules**

#### 4. Existing Approximate Adder Designs

Adders have become a performance bottleneck in the Arithmetic Logic Unit of DSP and general-purpose processors. Conventional design methodologies aimed at improving the performance of an adder frequently result in more circuit area and power consumption.

Approximations have become essential in adders and are implemented at various abstraction levels. Approximate adders compromise accuracy for an overall efficiency enhancement. Numerous approximate adder designs have been developed by reducing worst-case carry-chains, critical paths and hardware complexity. These approximate adder types include speculative adders, segmentation-based adders, Carry Select Adders

and Approximate Full Adders that are explained in the followingsubsections.

Table 3.1 Truth table of existing approximate FAs

<sup>a</sup><sub>×</sub> -Represents the inaccurate output

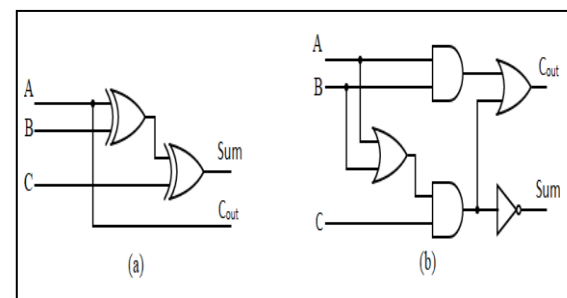
Inputs											
						InXA A1		InXA A2		InXA A3	
A	B	C	S	Cou t	S	C o ut	S	Cou t	S	C o ut	
0	0	0	0	0	0	0	0	0	1	×	0
0	0	1	1	0	1	1	×	1	0	1	0
0	1	0	1	0	1	0	1	0	1	1	0
0	1	1	0	1	0	1	1	×	1	0	1
1	0	0	1	0	1	0		0			0
1	0	1	0	1	0	1		1			1
1	1	0	0	1	0	0		1			1

Table 3.2. conventional Full Adder.

						SAF	SAF	SAF	SAF	
						A1E	A2E	A3E	A4E	
A	B	C	S	C <sub>out</sub>	S	C	S	C <sub>o</sub>	S	C <sub>o</sub>
0	0	0	0	0	0	0 <sup>out</sup>	0	0 <sup>ut</sup>	0	0 <sup>ut</sup>
0	0	1	1	0	1	0	1	0	1	0
0	1	0	1	0	1	0	1	0	1	0
0	1	1	0	1	0	1	1	0	1	0
							×	×	×	×
1	0	0	1	0	1	0	0	1	0	1
							×	×	×	×
1	0	1	0	1	0	1	0	1	1	1
									×	
1	1	0	0	1	0	1	0	1	1	1
									×	
1	1	1	1	1	0	1	1	1	0	1
					×			×		

## 5. PROPOSED SIMPLIFIED APPROXIMATE ADDERS(SAAs)

Based on the logic complexity reduction seven SAAs have been proposed by simplifying the K-Map of the Exact FA and HA designs. Circuits developed using this method have predictable errors, the designer can fix the number of logic gates based on the accuracy requirements set for a specific application.

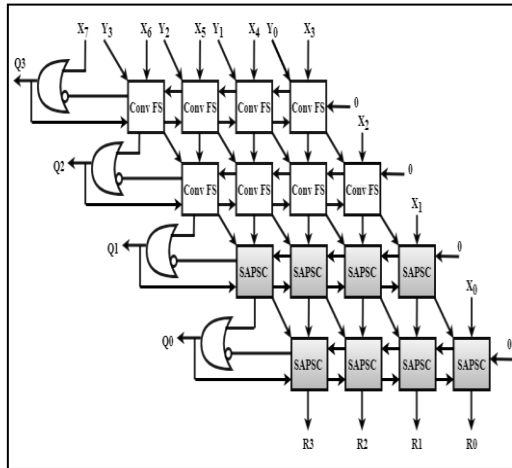


**Fig 7. Proposed structure**

SAAs are developed with novel logical expressions interpreted from its truth table to generate the Sum (S) and the Carry (Cout) to yield efficient area- delay optimization. In the proposed SAAs, (SAA1-SAA6) represents the approximate FA designs and SAA7 corresponds to an approximate HA design. The inputs of the following proposed SAAs are denoted by A, B and C, while the outputs are denoted by Sum and Cout. SAA3, SAA5 and SAA6 are designed to have errors in both the Sum and Cout.

while SAA1, SAA2 and SAA4 are designed to have errors either in Sum or output Cout.



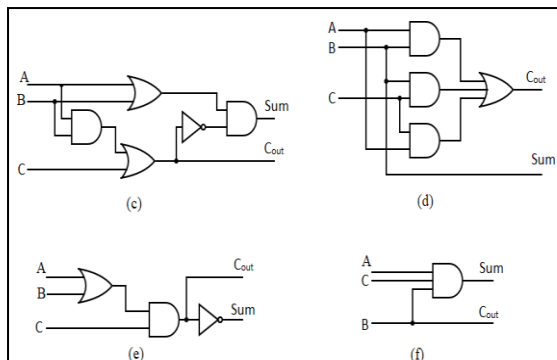


**Fig 8. Proposed structure implement**

$$Sum = (A + B)C_{out}'$$

$$C_{out} = AB + C$$

Gate level architecture of (c) SAA3 (d)



SAA4 (e) SAA5 (f) SAA6

## CONCULSION

Considering the extreme For extensive performance evaluation, both the proposed and existing designs are coded in Verilog Hardware Description Language (VHDL) and executed using Cadence Genus Synthesis Solution at 45 nm technology node. Various design metrics such as power, delay and area are considered for analysis. Considering approximate dividers, the proposed ARDXTRs and ARDXHRs exhibit improved performance than the existing designs. ARDXHR2 exclusively achieves a significant reduction in the PDP and in

comparison, with the exact design it achieves about 43.98% savings in power.

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