

# FINDING EIGEN VALUES FOR BUTTERFLY GRAPHS

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**Abstract** - In the mathematical field of graph theory the butterfly graphs (also called bowtie graphs and hourglass graphs) are planar undirected graphs with 5 vertices and 6 edges. The  $n$  - dimensional butterfly graphs, denoted by  $BF(n)$  has vertex set  $v = \{(x; i): x \in V(Q_n), 0 \leq i \leq n\}$ . Two vertices  $(x; i)$  and  $(y; j)$  are liked by an edge in  $BF(n)$  if and only if  $j = i+1$  and either  $x = y$ , or  $x$  differ from  $y$  in precisely the  $j^{\text{th}}$  bit for  $x = y$ , the edge is said to be straight edge, otherwise the edge is a cross edge, for fixed  $i$ , the vertex  $(x, i)$  is a vertex on level  $i$ . the  $n$ - dimensional butterfly graphs has  $2^{n(n+1)}$  vertices and  $2^{n+1}n$  edges. Butterfly graphs and domination are very important ideas in computer architecture and communication technique. In this paper, we present results about Eigen values of butterfly graphs  $BF(n)$ . We also show the Eigen value through adjacency matrix of butterfly graphs  $BF(n)$  are related to each other as  $a(BF(n)) = (BF(n)) + 1$ . If  $n = 2, 3, 4$  k.a  $(BF(n)) = (BF(n))$  if  $n = 4k$ .

**Key Words:** Graphs, planar graphs, butterfly graphs, butterfly network, Edge, Vertex, Eigen values.

## 1. INTRODUCTION

Graph theory itself is typically dated as beginning with Leonhard Euler's 1736 work on the seven bridges of Konigsberge. However drawing of another most popular bounded degree derivate network of the hypercube is what is called a butterfly network.

Butterfly network have been extensively studied in the literature because of their many application in computer architectures. In the undirected case, aspect that have been considered include, for example the study of the cycle structure and hamiltonicity[1,4,11] and the development of communication and routing algorithms[10,12,6-8]. Some of these results have been recently extended also the directed case[2,3]. The knowledge of this spectrum will facilitate a new approach for a better characterization of this important category of network. Throughout this paper we consider finite, simple and undirected graph  $G = (V(G), E(G))$  with  $p$  vertices and  $q$  edges.  $G$  is also called a  $(p, q)$  graph. We follow the basic notations and terminologies of graph theory as in [5].

## 2. EIGENVALUES OF GRAPHS

The Eigen value problems appear in various contexts throughout mathematics and engineering and refer to determining all possible list of Eigen value(spectra) for matrices fitting some description. Graphs often describe relationship in a physical setting, such as control of a system and the Eigen value of associated matrices govern the behavior of the system. The Eigen value problem of a given graph  $G$  is to determine all possible spectra of real symmetric matrices whose off-diagonal entries are governed by the adjacencies in  $G$ . It was thought by many researchers in the field that atleast for a tree  $T$ , determining the ordered multiplicity list of  $T$  would suffices to determine the spectra of matrices describe by  $T$ . when it was show in [9] that this was not the case, the focus of much of the research in the area shifted to the narrower question of maximum Eigen value multiplicity or equivalently maximum nullity or minimum rank of matrices described by the graphs . The Eigen values of the graphs are closely tied to the Eigen values of a graph. while you do not always need to find the Eigen values to find Eigen values in some cases using the structures of a graphs we can construct the eigenvector and thus find the Eigen values.

A famous question in spectral graph theory is "can you hear the shape of a graphs? That is given the Eigen values can you determine the graphs that produced them.

## SPECTRAL GRAPH THEORY

One of the most useful ways of doing this has been by studying the various spectra of matrices ( that is the Eigen values of the matrices) that can be associated with the graph. By facing at these Eigen values it is possible to get information about the graphs that might otherwise be difficult to obtain the study of the relation between Eigen values is the heart of spectral graphs theory thus someone interested in using spectral graphs theory needs to be familiar both with graph theory and the basic tools of liner algebra including Eigen values ,Eigen vectors, determinants , the courant-fischer theorem, the perron-frobenius theorem and so on...

we will introduce the most common matrices associated with graphs namely the adjacency matrix and given some simple example for each about how the Eigen values can be used to give some information about the graphs.

### THE ADJACENCY MATRIX

Given a graph G we can form a matrix A called the adjacency matrix, by letting the vertices index the columns and rows and then letting

$$A_{ij} = \begin{cases} 1 & \text{if } i \text{ is adjacent to } j \\ 0 & \text{if } i \text{ is not adjacent to } j \end{cases}$$

Let G be a simple graph of order n and let A(G) be its adjacency matrix. The Eigen values  $\lambda_1, \lambda_2, \dots, \lambda_n$  of A(G) are referred to as the Eigen

values of the graph G and form the spectrum of G [19][7].

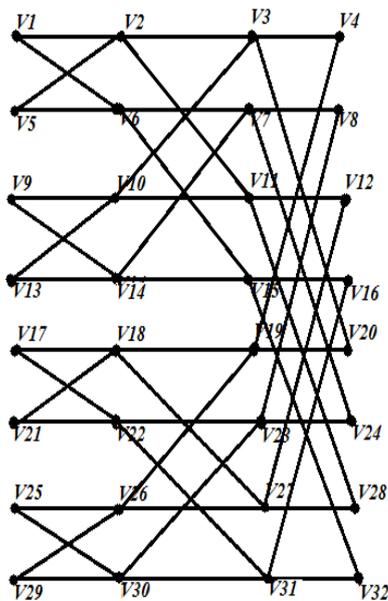
### BUTTERFLY NETWORK

The **n-dimensional** butterfly network denoted by BF(n), has vertex set  $V = \{(x;i) : X \in V(Q_n), 0 \leq i < n\}$ .

The vertices (x;i) and (y;j) are linked by an edge in BF(n) if and only if  $j = i + 1$  and either

- i)  $X = Y$
- ii) X differ from y in precise the  $j^{th}$  bit.

For  $x = y$ , the edge is said to be straight edge, otherwise the edge is a cross edge for fixed I, the vertex (x,i) is vertex on level i. the graphs is show in BF(3)



BF(3)

The **n-dimensional** butterfly graphs has  $2^n(n+1)$  vertices and  $2^{n+1}.n$  edges.

The butterfly network BF(n) has  $2^n(n+1)$  vertices because BF(n) has n+1 levels and there are  $2^n$  vertices in every level. each vertex on level 0 and n is of degree 2, otherwise, every is of degree 4. It is clear that BF(n) is eulerian since it is connected and has no odd degree vertices[6].

### BUTTERFLY NETWORK BF(r)

The definition of the butterfly network BF(r) is as follows: The r-dimensional butterfly has  $(r+1)2^r$  nodes and  $r.2^{r+1}$  edges. The set V of nodes of an r-dimensional butterfly correspond to pairs [w,i], where i is the level of a node ( $0 \leq i \leq r$ ) and w is an r-bit binary number that denotes the row of the node. Two nodes  $\langle w, i \rangle$  and  $\langle w_0, i_0 \rangle$  are linked by an edge if and only if  $i_0 = i + 1$  and either w and  $w_0$  are identical, or w and  $w_0$  differ in precisely the ith bit. An r-dimensional butterfly is denoted by BF(r).

**Theorem** : Let G be an r-dimensional butterfly network.

$$\text{Then } \beta(G) = \left\lceil \frac{r+1}{2} \right\rceil . 2^r$$

**Proof** : Let G be an r-dimensional butterfly network.

**Case(i)**: If r is even then the number of levels which is r + 1 is an odd number. Therefore we require exactly  $\left\lceil \frac{r+1}{2} \right\rceil . 2^r$  vertices to cover all edges.

**Case(ii)**: If r is odd then the number of levels which is r + 1 is an even number. So we need exactly  $\left\lceil \frac{r+1}{2} \right\rceil . 2^r$  vertices to cover all the edges.

$$\text{Combining both the cases, } \beta(G) = \left\lceil \frac{r+1}{2} \right\rceil . 2^r$$

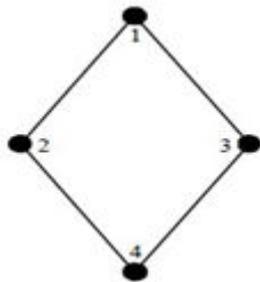
**Corollary** : Let G be an r-dimensional butterfly network. Then (i) If r is odd, then  $\beta(G) = \frac{|V|}{2}$ . (ii) If r is even, then  $\beta(G) < \frac{|V|}{2}$ .

**Theorem**: Let G be an r-dimensional butterfly network. Then  $\beta(\overline{G}) = (r+1).2^r - 2$ .

**EIGEN VALUE WITH BUTTERFLY GRAPHS**

**BUTTERFLY GRAPH(1)**

Adjacent matrix = 
$$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$



$$\det(\lambda I - A) = \det \begin{bmatrix} \lambda & -1 & -1 & 0 \\ -1 & \lambda & 0 & -1 \\ -1 & 0 & \lambda & -1 \\ 0 & -1 & -1 & \lambda \end{bmatrix}$$

$$\det(\lambda I - A) = \lambda \begin{bmatrix} \lambda & 0 & 0 & 0 \\ 0 & \lambda & 0 & 0 \\ 0 & 0 & \lambda & 0 \\ 0 & 0 & 0 & \lambda \end{bmatrix} - \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

$$\det(\lambda I - A) = \lambda \begin{bmatrix} \lambda & 0 & -1 \\ 0 & \lambda & -1 \\ -1 & -1 & \lambda \end{bmatrix} - 1 \begin{bmatrix} -1 & 0 & -1 \\ 0 & -1 & \lambda \end{bmatrix} - 1 \begin{bmatrix} -1 & \lambda & -1 \\ -1 & 0 & -1 \end{bmatrix} + 0 \begin{bmatrix} -1 & \lambda & 0 \\ -1 & 0 & \lambda \end{bmatrix}$$

$$\det(\lambda I - A) = \lambda [\lambda(\lambda+1) - 1(0-\lambda)] - 1[-1(\lambda+1) - 1(1)] - 1[-1(1) - \lambda(-\lambda-0) - 1(1-0)]$$

$$\det(\lambda I - A) = \lambda[\lambda^2 + \lambda + \lambda] - 1[-\lambda - 1 - 1] - 1[-1 + \lambda^2 - 1]$$

$$\det(\lambda I - A) = \lambda[\lambda^2 + \lambda - \lambda + \lambda + 1 - 0 + 1 - 1 + \lambda^2 + 1]$$

$$\det(\lambda I - A) = \lambda[2\lambda^2 + \lambda + 2]$$

$$\lambda = -2, 0, 0, 2$$

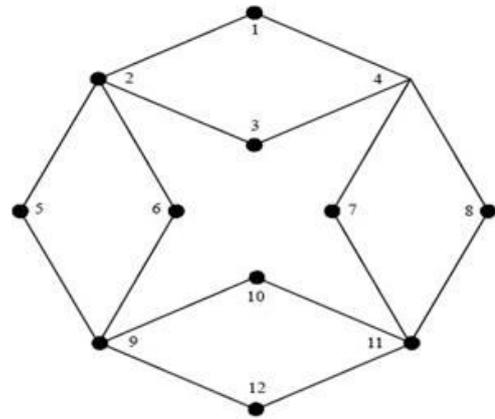
**Eigen value  $\lambda$ : -2, 0, 0, 2**

**Maximum value: 2**

Using matlab we can find Eigen values for **n vertex**

Now we can find Eigen value for **4 vertex**

**BUTTERFLY GRAPH(2)**



Adjacent matrix =

$$\begin{bmatrix} 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

**Eigen value  $\lambda$ : 2, -2, -2, 0, 0, 0, 0, 0, 0, 2, 2, 2**

**Maximum value  $\lambda$ : 2**

We can find Eigen values for **12 vertex**

Eigen values for **12 vertex using matlab**



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## BIOGRAPHY



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