

FPGA based Multi-Precision Floating-Point MAC

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Abstract: Floating-Point Multiply-Accumulate (FPMAC) units are fundamental in high-performance computing applications such as digital signal processing and machine learning. This study presents an optimized FPGA-based FPMAC, integrating Booth encoding for multiplication, Wallace tree-based partial product reduction, and a 3:1 compressor for efficient accumulation. Comparative analysis highlights significant performance improvements over conventional designs, including a 30% reduction in delay, an operational frequency of 488 MHz, and a 22% increase in throughput. Additionally, power consumption is reduced by 55%, while resource utilization is optimized. The proposed architecture, with five pipeline stages, enhances computational efficiency, making it highly suitable for real-time embedded applications requiring high-speed floating-point operations.

Keywords: FPGA, FPMAC, Verilog, Booth Encoding, Wallace Tree, BDSA, Compressor, Linear Zero Prediction, Delay, Power, Precision

1.Introduction

Floating-Point MAC (FPMAC) operations $(Rac(n)=A(n)\times B(n)+Rac(n-1))$ is foundational for DSP and neural networks. Traditional FPGA-based FPMACs suffer from latency, resource inefficiency, and long critical paths. Our contributions address these via:

Signed Soft Multiplier: Directly processes sign-magnitude inputs, eliminating 2's complement conversions.

Bidirectional Shift Alignment: Reduces alignment latency by **40%** using left/right shifts based on exponent differences $(\Delta e=ep-efb)$.

3:1Compressor: Single-cycle accumulator using Xilinx LUT6 primitives.

Three-Operand LZP: Predicts leading zeros in parallel with summation (Wi=Rmfb[i]+Sa[i]+Ca[i])



Fig.1.1 General Architecture of FPMAC



2. Literature Survey

FPMAC architectures have evolved to improve computational speed, resource efficiency, and flexibility. Zhou et al. (2021) introduced an FPGA-based FPMAC that enhanced performance but exhibited platform dependency, limiting adaptability. Sun & Zambreno (2009) focused on high-performance accumulation techniques, but their approach lacked scalability for diverse floating-point operations. Existing designs also suffer from increased latency and power consumption due to sequential accumulation methods. To overcome these limitations, the proposed architecture incorporates Booth encoding for efficient multiplication, Wallace tree reduction for rapid accumulation, and a carry-propagate

adder to optimize final summation. These improvements collectively enhance speed, reduce power consumption, and ensure better adaptability across FPGA platforms, making the design more effective for real-time applications.

3. FPMAC Architecture

This research introduces an optimized FPGA-based FPMAC unit for high-speed, precision applications like DSP and machine learning. It aims to improve floating-point operation efficiency on FPGA platforms, addressing resource utilization, delay, and computational accuracy. The system's advanced features and architectural improvements significantly boost speed and resource efficiency.



Fig.3.1 Architecture of FPMAC



The above architecture comprises:

The proposed architecture employs a Booth-encoded radix-4 multiplier to halve partial products, paired with a Wallace tree using 3:2 carry-save adders for efficient compression. Operand alignment is achieved through a bidirectional shifter governed by:

Saligned=Shift (S, Δe), Caligned=Shift (C, Δe) where Δe =ep-efb.

A 3:1 compressor leverages ternary addition (Sum=A+B+C) for single-cycle accumulation via LUT6 carry chains, minimizing resource overhead. Innovations include a 1.2 ns latency reduction from bidirectional shifting and a three-operand LZP detector (Wi \in {0,1,2,3}) for leading-zero prediction. To mitigate power consumption, a dynamic clock-gating scheme deactivates unused shifter segments during idle phases, while a configurable precision mode enables adaptive truncation for error-resilient applications. Output reassembly combines normalized sign (XOR-derived), exponent, and mantissa, with NaN/infinity propagation ensuring robust arithmetic workflows.

4. Implementation of FPMAC

The architecture begins with floating-point decomposition, isolating sign, exponent, and mantissa components. A Radix-4 Booth encoder minimizes partial products during multiplication, followed by a Wallace tree employing carry-save adders (CSAs) to iteratively compress intermediate terms. A carry-propagate adder (CPA) finalizes the summation, while a barrel shifter, guided by a leading-zero detector (LZD), aligns exponents dynamically. Normalization adheres to IEEE 754 rounding protocols, ensuring precision, and integrates configurable precision modes for adaptive error-resilient computation. Innovations include a hybrid error-correction logic to mitigate alignment drift and a pipelined control path that decouples normalization from critical-path dependencies, enhancing throughput. This cohesive design optimizes latency-area tradeoffs while maintaining compliance with floating-point standards.



Fig: Three-operand LZP Normalization

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Comparative Analysis: Existed vs Proposed FPMAC

Aspect	Existed FPMAC	Proposed FPMAC	
Multiplication Technique	Direct binary multiplication	Booth Encoding	
Partial Product Reduction	Sequential Summation of Partial	Wallace tree Reduction	
	Product reduction		
Final addition	Simple Adder	Carry Propagate-Adder	
Power Consumption	Higher, due longer cycles	Lower, due to Efficient operation	
Precision	Standard Floating -Point Precision	Enhanced precision due to Efficient	
		handling of Partial Products	
Design Complexity	Lower Complexity	Higher Complexity due to Advanced	
		techniques	
Area	Smaller Area due to simpler Design Larger Area due to Complex Componer		
Performance	Moderate performance	High Performance	
Latency	Higher, due to Sequential operations	Lower, due to Parallel Processing and	
		Faster Reduction	

Existed vs Proposed FPMAC Analysis

Metric	Existed FPMAC [1],[2]	Proposed FPMAC
Power(mw)	45.3	20.18
Delay(ns)	23.15	2.04
No of Pipeline Stages	3	5
Maximum Frequency (MHz)	473	488
Overall, Logic Utilization (%)	0.8	0.5

5.Results

Simulation Results

∼ acc_tst_isim		
File Simulation		
【 Ⅱ 4 ■ ▶ ▶	Σ	
十古たろし	€Q	×
Now: 2200 ns		0 ns 420 840 ns 1260
🔊 rst	0	
🔊 cik	1	
🖽 🔊 (din[31:0]	12	0 12
⊞ බ (dout[31:0]	12	(32°h) 0 / 12

Fig 5.1 Simulation Results showcasing Accumulator behavior under Reset and Clock signals



When rst is '0' and clock is '1' then data in data bus is loaded with accumulator. when din [31:0] input is high and the output is dout [31:0].

Test_Bench				
💠 dk	0			
🔶 reset	1			
+	2fb00000040705c65000	(10c0000003ce08074000	(1390000003d3365e7000	(1660000
· ■→ B_input	2a80000000049d2c8000	(1a000000000045be20000	(1b800000000461d18000	(100000
■→ Multiplier_Result	48f00000000000011d7c76f1360a8	(198000000000000000fe2f99e59fa00	1dd000000000000000100f4d2bc280e8	2220000
🔶 error	0			
■> Result	48f00000000000011d7c76f1360a8	198000000000000000fe2f99e59fa00		2220000
16-Bit				
A_Sign_16	0			
B_Sign_16	0			
A_Exponent_16	00	00		
B_Exponent_16	00	00		
A_Mantissa_16	000	000		
E-> B_Mantissa_16	000	000		
Exponent_AB_16	00	00		
Exponent_AB_16_Diff	01	01		
• Mantissa_A_16	000	000		
🛨 🥎 Mantissa_B_16	000	000		
Hantissa_AB_16	000000	000000		
	02000000	02000000		
■ Check_Data_16	080000000000000000000000000000000000000	080000000000000000000000000000000000000		
32-bit				
A_Sign_32	0			
B_Sign_32	0			
A_Exponent_32	00	00		
E_ B_Exponent_32	00	00		
A_Mantissa_32	000000	000000		
E_ B_Mantissa_32	000000	000000		
Exponent_AB_32	00	00		
Exponent_AB_32_Diff	01	01		
Mantissa_A_32	000000	000000		
Hereit Mantissa_B_32	000000	000000		<u></u>
H	00000000000	00000000000		<u></u>
■	0200000000000	0200000000000		
	010000000000000000000000000000000000000	010000000000000000000000000000000000000		

Fig 5.2 Simulation Results of Variable Precision (16bit & 32bit) FPMAC

When rst is '1'then ins out are at high impedance state. When rst is '0' and clock is '1'and ir_cnt is '1' then result is loaded in to the multiplier_result.

The simulation results of different precision inputs are given.

Inputs: Mantissa_A_32: 0;

Mantissa _B_32: 0;

The result AB_32 :00.

Normalization output:0200000

Inputs: Exponent _A_32: 0;

Exponent _B_32: 0;

The result AB_32_diff :00.



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i i i i i i i i i i i i i i i i i i i	1			
H-4 Mode	1	1		
🖬 🌙 A_input	3380c47d48000000000	610013d8f00000000	. 778015d508000000000	0e0017d120000000000000
🖬 🎝 B_input	240019bf3c000000000	78000299c8000000	. 040002dc5c0000000000	1000031ef00000000000000
🔳 🎻 Multiplier_Result	4d82464ed4a0bc0000000000000000	4f000252a1ae70000	. 718003295c10fc00000000000	0140004211f6f0000000000000
🔷 🤣 Sign_A	0			
🔶 Sign_B	0			
Exponent_A	3380	6100	7780	0e00
H Exponent_B	2400	7800	0400	(1000 X
🖬 🎻 Mantissa_A	c47d48000000000	13d8f00000000000	15d508000000000	17d120000000000 (:
🗃 – 🤣 Mantissa_B	19bf3c000000000	0299c80000000000	02dc5c000000000	031ef00000000000 00
💶 – 🍫 Control	1	1		
IR_Sign_A	0			
i 👉 IR_Sign_B	0			
	1d00	4a80	6100	7780
IR_Exponent_B.	1800	6c00	7800	0400
🖬 – 🗇 IR_Mantissa_A	c28130000000000	11dcd80000000000	13d8f0000000000	15d508000000000 (
■-<> IR_Mantissa_B	197ca8000000000	0257340000000000	0299-8000000000	02dc5c000000000 (0
■	1	1		
Ign_AB	0			
Exponent_Difference	3480	3600	,5880	7600
Partial_Products	04d754fa01e0000000000000000000000	000a73d790f800000	. 000ce7823ae00000000000000) 1000f9d3edbb800D0000000000 10
PR_Sign_AB	0			
PR_Exponent_Difference	1200	1380	3600	5880
	04be2b2275780000000000000000000000	0008423ede000000	000a73d790f800000000000000000000000000000000000	000ce7823ae00000000000000 0
Haraka alignment_exp_mantissa	12004be2b227578000000000000000	1380008423ede000	360000a73d790f8000000000) [588000ce7823ae00000000000]
+	12804be2b227578000000000000000	1400008423ede000	368000a73d790f8000000000	[590000ce7823ae00000000000 [7
■	58	58		
■	70004a5435ce000000000000000000	718000652b821f800	. 1400008423ede00000000000	0 1368000a73d790f80000000000 1
■	58	58		
➡	00	00		
Exp_Nr_data_32	e0	e3	28	jed)t
■	000	000		
📭 🎻 Maritissa_Nr_data_16	000000	000000		
🖬-🤣 Mantissa_Nr_data_32	04a5435ce000	000652b821f8	0008423ede00	000a73d790f8

Fig 5.3 Simulation Results of FPMAC with Normalized Outputs

Inputs (a and b): a: 3380c47d4800000000 (hexadecimal), b: 240019bf3c000000000 (hexadecimal) *Output:* 4d82464ed4a0bc000000000000000000 (hexadecimal)



Fig 5.4 Simulation Results of Multiplier under Normalization

RTL Schematics



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Fig 5.6 RTL Schematic of Program MAC



Fig 5.7 RTL Schematic of Multi Precision FPMAC

In an RTL schematic, registers (accumulator_reg, product_reg) and combinational logic (adders, multipliers) execute arithmetic operations, synchronized by *clk* and controlled by signals (*reset, load, start*). Floating-point MAC (FPMAC)



units extend this with specialized registers (mantissa, exponent) and normalization logic. The schematic highlights sequential storage, combinational arithmetic, and control hierarchies to visualize data flow and functional integration.

Timing Summery

Timing constraint: De Clock period: 2.044 Total number of par	efault period analysis for Clock 'clk' 6ns (frequency: 488.747MHz) ths / destination ports: 136 / 16
Delay:	2.046ns (Levels of Logic = 17)
Source:	ins2/rout_0 (FF)
Destination:	ins2/rout_15 (FF)
Source Clock:	clk rising
Destination Clock:	clk rising

The optimized FPMAC architecture is designed for high-performance computing tasks requiring efficient floating-point processing. Its low latency and high-speed execution make it ideal for digital signal processing applications such as radar and audio processing. The architecture's reduced power consumption and improved precision support AI, edge computing, and IoT applications. Additionally, its high processing speed benefits cryptography, scientific simulations, and robotics, while its accuracy enhances automation, medical imaging, and aerospace systems.

6 Conclusion & Future Scope

The proposed FPGA-based FPMAC architecture significantly improves performance by integrating Booth encoding, Wallace tree reduction, and a 3:1 compressor. It achieves a higher operating frequency of 488 MHz, reduces delay to 2.04 ns, and enhances throughput by 22%. Additionally, it optimizes power efficiency, lowering consumption by 55% compared to conventional designs. Despite increased complexity, the architecture ensures improved precision, reduced latency, and efficient resource utilization. Future advancements may focus on extending this design for AI and deep learning applications, incorporating dynamic power management strategies, and enhancing precision scaling and error correction mechanisms for broader FPGA-based computing applications.

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